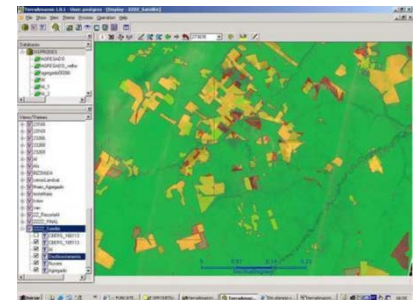
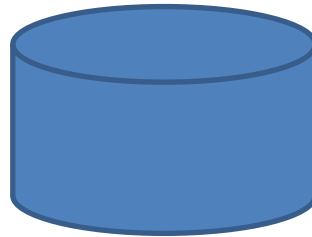
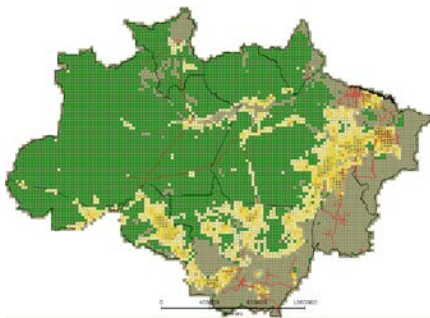


# Spatial Databases: Lecture 4

Institute for Geoinformatics  
Winter Semester 2014



Malumbo Chipofya: room 109

# Topic Overview

1. Prelude: Data and problem solving in science and applications
- 2. The Relational Database model**
- 3. Interacting with relational databases**
4. Spatial Relational Database Management Systems
5. Applications: Terraview and Terralib: Prof. Dr. Gilberto Camara
6. A sample of Nosql Databases: brief introductions + example applications
  - a. Array databases: SciDB
  - b. Document databases: MongoDB
  - c. Graph databases: Neo4J
7. Summary of all lectures given.

# Recap

- Candidate Keys:
  - Uniqueness + Irreducibility
- Relational Operations:
  - Restrict + Project + Join
- Functional Dependence:  $B \rightarrow A$ 
  - A is functionally dependent on B
  - B functionally determines A

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# Functional Dependencies

- Given two sets of attributes of a relation  $R$ :

$$A := \{a, b, c, \dots\} \quad B := \{x, y, z, \dots\}$$

- $A$  is a functionally dependent on  $B$  written

$$B \rightarrow A$$

***if and only if*** there is a function from the set of legal values of  $B$  to the set of legal values of  $A$  determined exactly by tuples of  $R$

# Functional Dependencies

- Trivial FD
  - $LHS \supseteq RHS$
- The closure of a set  $S$  of FDs (denoted  $S^+$ )
  - The set of **all** FDs that can be derived from  $S$
  - $S^+$  can be computed using few simple rules



# Functional Dependencies

- Rules – we write 'A' for {A} and 'A,B,C' for {A, B, C}
  - Reflexivity:
    - $B \subseteq A$  implies  $A \rightarrow B$
  - Augmentation:
    - $A \rightarrow B$  implies  $A,C \rightarrow B,C$
  - Transitivity:
    - $A \rightarrow B$  and  $B \rightarrow C$  implies  $A \rightarrow C$
  - Self-determination:
    - $A \rightarrow A$
  - Decomposition:
    - $A \rightarrow B,C$  implies  $A \rightarrow B$  and  $A \rightarrow C$
  - Union:
    - $A \rightarrow B$  and  $A \rightarrow C$  implies  $A \rightarrow B,C$
  - Composition:
    - $A \rightarrow B$  and  $C \rightarrow D$  implies  $A,C \rightarrow B,C$

# Functional Dependencies Example:

$$\{A \rightarrow B, C; C \rightarrow D\}$$

- Reflexivity:  $B \subseteq A$  implies  $A \rightarrow B$
- Augmentation:  $A \rightarrow B$  implies  $A, C \rightarrow B, C$ 
  - $A, C \rightarrow A, D; A, C \rightarrow B, C; A, D \rightarrow B, C, D; B, C \rightarrow B, D$
- Transitivity:  $A \rightarrow B$  and  $B \rightarrow C$  implies  $A \rightarrow C$ 
  - $A \rightarrow B, D; A, C \rightarrow B, D; A, C \rightarrow B, C, D$
- Self-determination:  $A \rightarrow A$
- Decomposition:  $A \rightarrow B, C$  implies  $A \rightarrow B$  and  $A \rightarrow C$ 
  - $A \rightarrow B; A \rightarrow C; A \rightarrow D;$
- Union:  $A \rightarrow B$  and  $A \rightarrow C$  implies  $A \rightarrow B, C$ 
  - $A \rightarrow C, D;$
- Composition:  $A \rightarrow B$  and  $C \rightarrow D$  implies  $A, C \rightarrow B, D$ 
  - *Anything else?*

# Functional Dependencies

- Irreducibility
  - A set of FDs,  $S$ , is irreducible if and only if it satisfies
    - RHS of every FD in  $S$  has only one attribute
    - LHS of every FD in  $S$  is irreducible in the sense that discarding any attribute changes the closure of  $S$  – **left irreducibility**
    - Discarding any FD in  $S$  changes the closure of  $S$

# Functional Dependencies

- Irreducibility: from last example

–  $A, C \rightarrow A, D$

–  $A, C \rightarrow B, C$

–  $A, D \rightarrow B, C, D$

–  $B, C \rightarrow B, D$

–  $A, C \rightarrow B, D$

–  $A, C \rightarrow B, C, D$

# Functional Dependencies

- Irreducibility: from last example

- RHS of every FD in  $S$  has only one attribute  
(Decomposition)

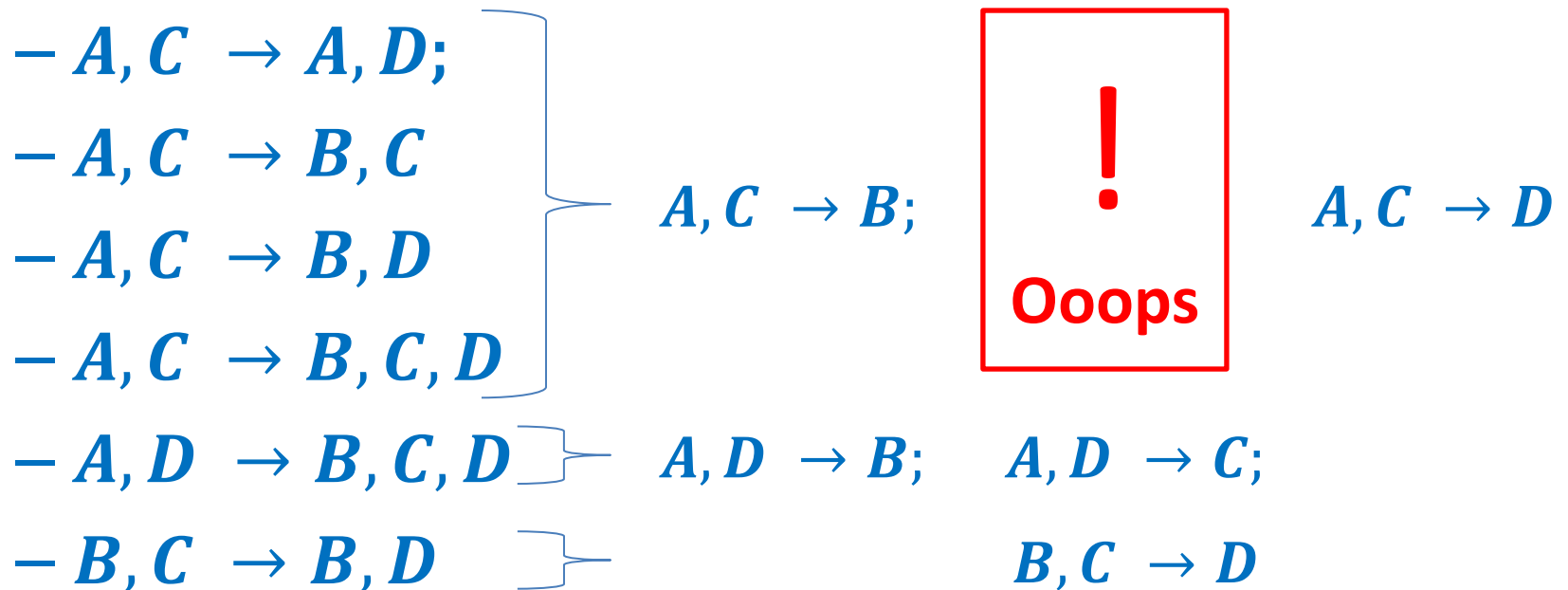
$- A, C \rightarrow A, D;$	}	$A, C \rightarrow B;$	$A, C \rightarrow B;$	$A, C \rightarrow D$
$- A, C \rightarrow B, C$				
$- A, C \rightarrow B, D$				
$- A, C \rightarrow B, C, D$				
$- A, D \rightarrow B, C, D$	}	$A, D \rightarrow B;$	$A, D \rightarrow C;$	$A, D \rightarrow D$
$- B, C \rightarrow B, D$	}	$B, C \rightarrow B;$	$B, C \rightarrow C$	

# Functional Dependencies

- Irreducibility: from last example
    - Discarding any FD in  $S$  changes the closure of  $S$   
(Discard the trivial FDs + all those that can be derived)
    - $A, C \rightarrow A, D;$
    - $A, C \rightarrow B, C$
    - $A, C \rightarrow B, D$
    - $A, C \rightarrow B, C, D$
    - $A, D \rightarrow B, C, D$
    - $B, C \rightarrow B, D$
- $A, C \rightarrow B;$      $A, C \rightarrow C;$      $A, C \rightarrow D$
- $A, D \rightarrow B;$      $A, D \rightarrow C;$      $A, D \rightarrow D$
- $B, C \rightarrow B;$      $B, C \rightarrow D$

# Functional Dependencies

- Irreducibility: from last example
  - Discarding any attribute on LHS changes the closure of S – **left irreducibility**



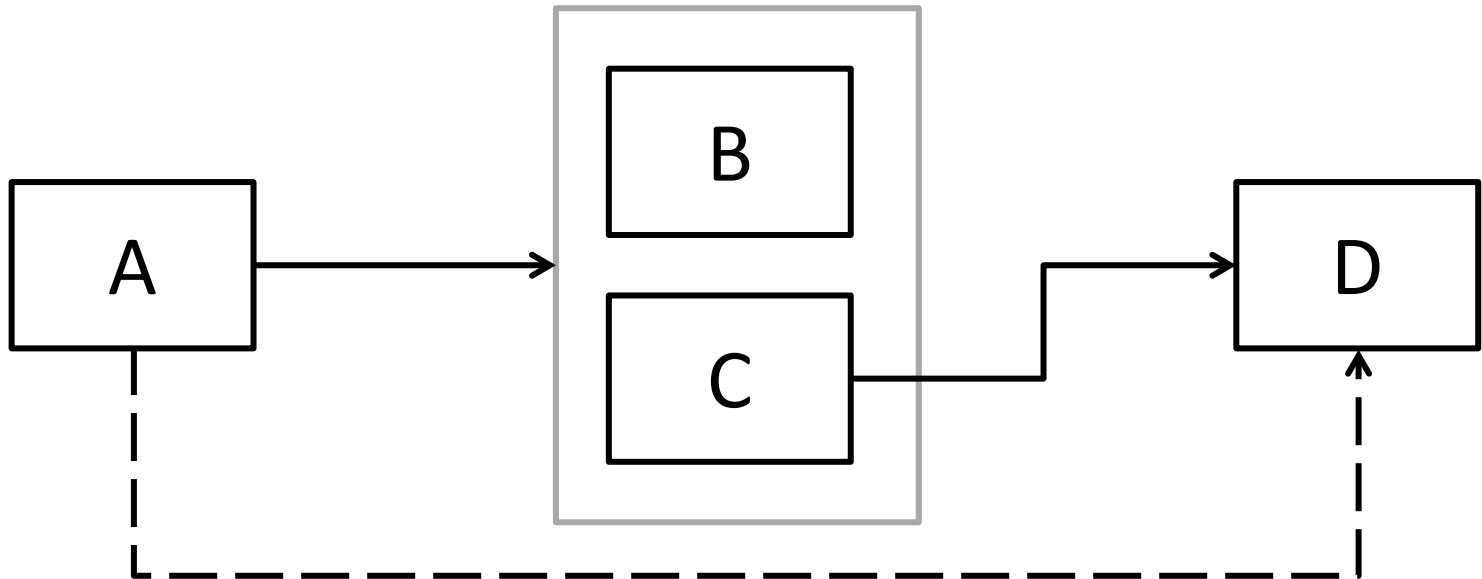
# Functional Dependencies

- Irreducibility: from last example
  - $A \rightarrow B, C$
  - $C \rightarrow D$ 
    1.  $A \rightarrow B$
    2.  $A \rightarrow C$
    3.  $C \rightarrow D$
- The irreducible equivalent is **NOT** unique



# Functional Dependency Diagrams

- $\{A \rightarrow B, C; C \rightarrow D\}$



# Normal Forms

- Example: Consider the our relation

ID#	Skill	M.St	#Chd	#Yrs	M.€	Date	#sticks	Wgt.	Hrs
1	Medium	M	0	2	40	1.06	55	9	6
2	Low	S	0	1	30	7.05	34	5	5
3	High	S	2	3	45	1.06	54	9	6
4	High	M	3	4	50	3.11	61	12	8

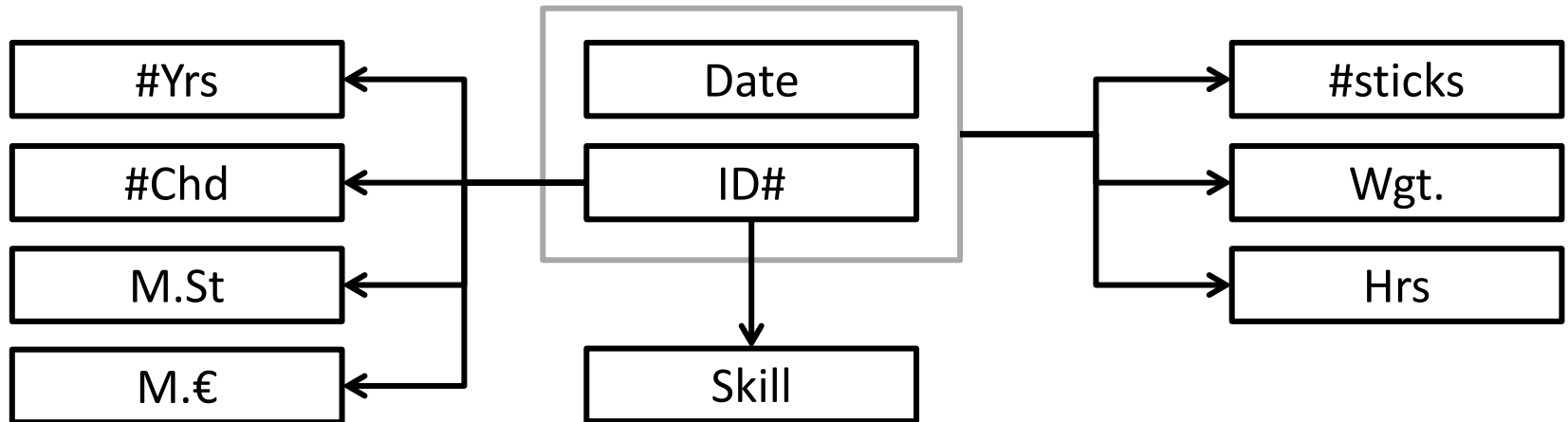
## 1<sup>st</sup> Normal Form (1NF)

- All legal relations are in 1NF

# Normal Forms

- Some FDs in this relation?

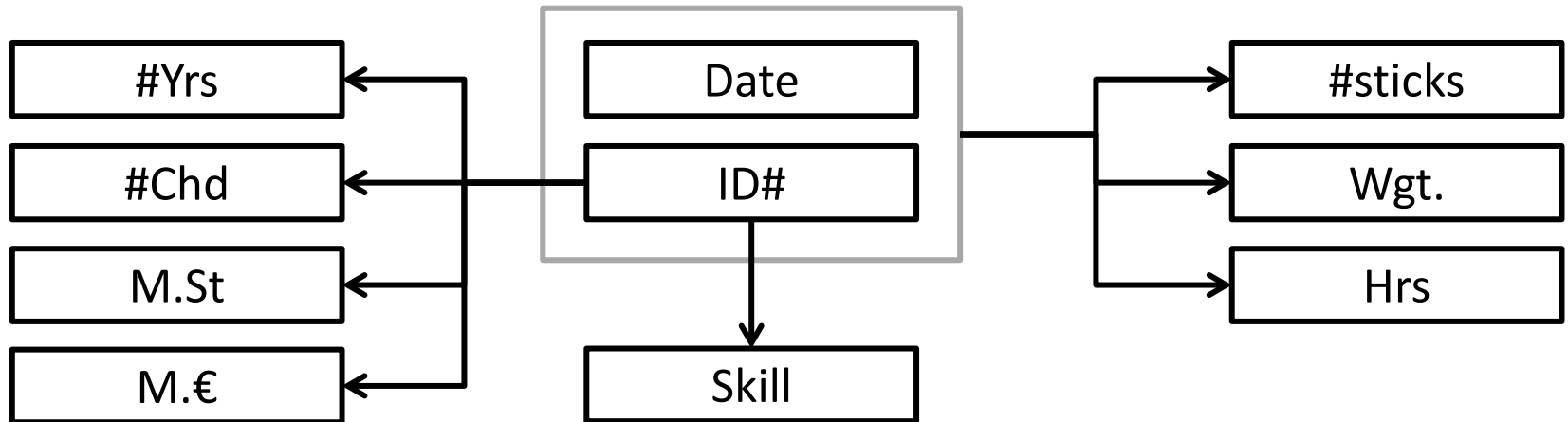
ID#	Skill	M.St	#Chd	#Yrs	M.€	Date	#sticks	Wgt.	Hrs
1	Medium	M	0	2	40	1.06	55	9	6
2	Low	S	0	1	30	7.05	34	5	5
3	High	S	2	3	45	1.06	54	9	6
4	High	M	3	4	50	3.11	61	12	8



# Normal Forms

- What are the problems with this relation?

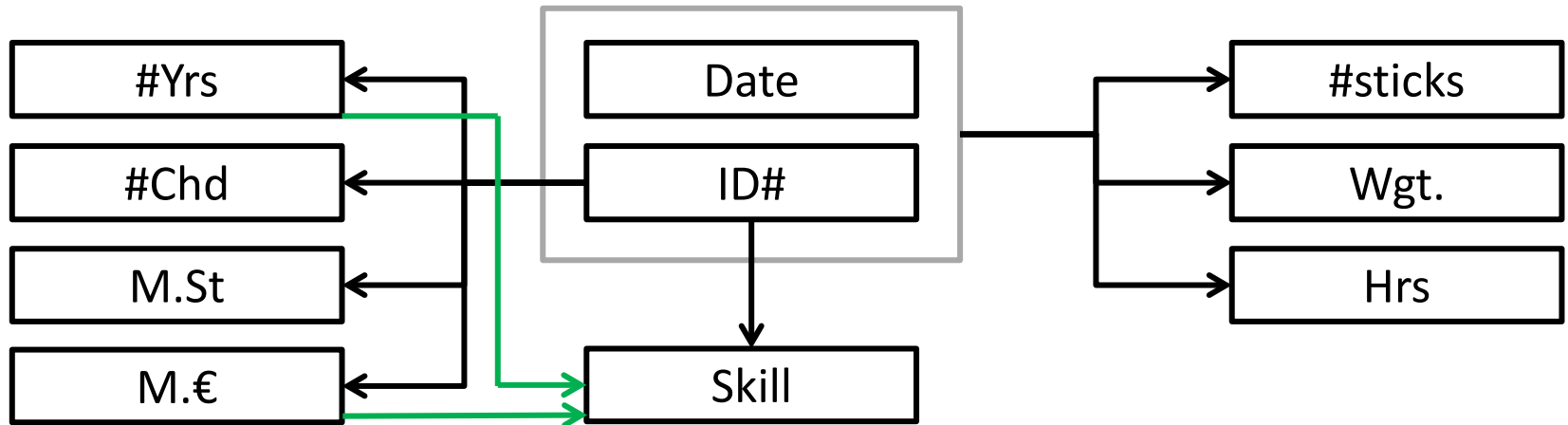
ID#	Skill	M.St	#Chd	#Yrs	M.€	Date	#sticks	Wgt.	Hrs
1	Medium	M	0	2	40	1.06	55	9	6
2	Low	S	0	1	30	7.05	34	5	5
3	High	S	2	3	45	1.06	54	9	6
4	High	M	3	4	50	3.11	61	12	8



# Normal Forms

- Let's reveal a few more dependencies

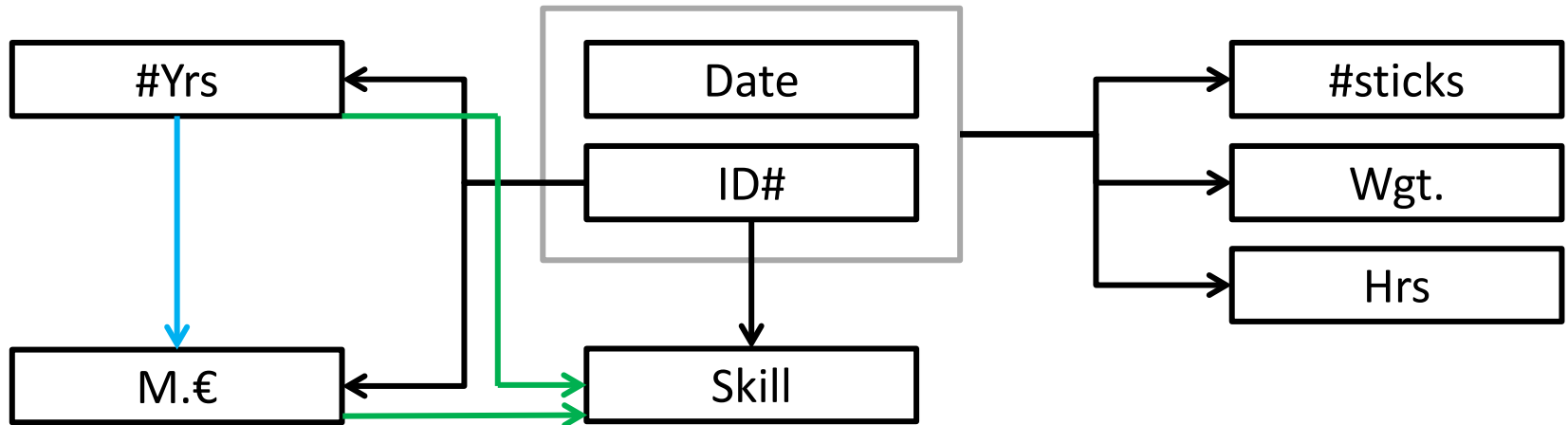
ID#	Skill	M.St	#Chd	#Yrs	M.€	Date	#sticks	Wgt.	Hrs
1	Medium	M	0	2	40	1.06	55	9	6
2	Low	S	0	1	30	7.05	34	5	5
3	High	S	2	3	45	1.06	54	9	6
4	High	M	3	4	50	3.11	61	12	8



# Normal Forms

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ID#	Skill	M.St	#Chd	#Yrs	M.€	Date	#sticks	Wgt.	Hrs
1	Medium	M	0	2	40	1.06	55	9	6
2	Low	S	0	1	30	7.05	34	5	5
3	High	S	2	3	45	1.06	54	9	6
4	High	M	3	4	50	3.11	61	12	8

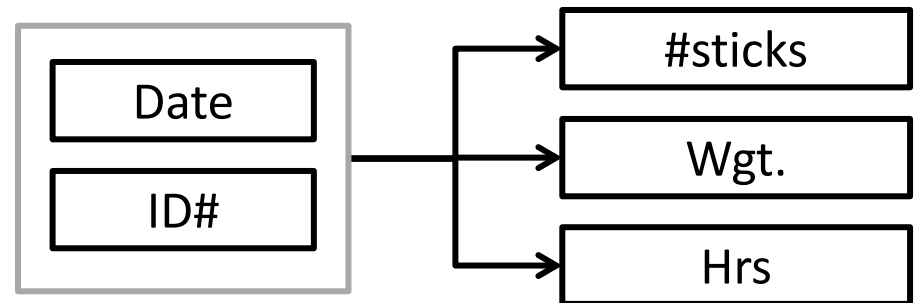
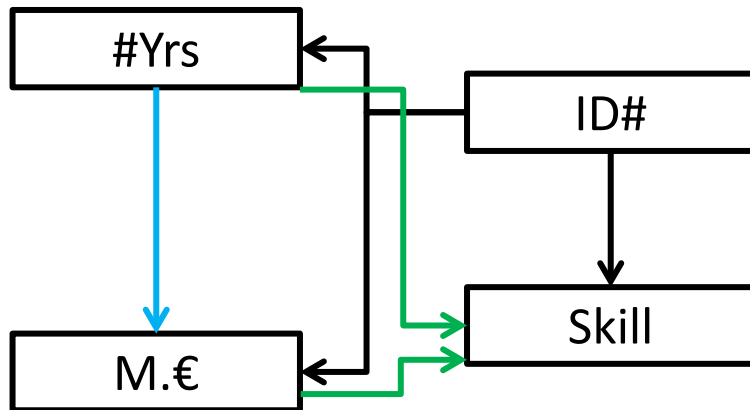


# Normal Forms

- Decompose the relation by projecting it

ID#	Skill	#Yrs	M.€
1	Medium	2	40
2	Low	1	30
3	High	3	45
4	High	4	50

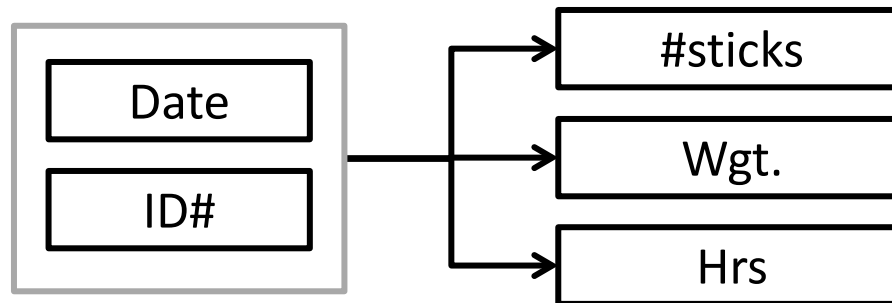
ID#	Date	#sticks	Wgt.	Hrs
1	1.06	55	9	6
2	7.05	34	5	5
3	1.06	54	9	6
4	3.11	61	12	8



# Normal Forms

- This relation is fine – It's at least in 2NF

ID#	Date	#sticks	Wgt.	Hrs
1	1.06	55	9	6
2	7.05	34	5	5
3	1.06	54	9	6
4	3.11	61	12	8

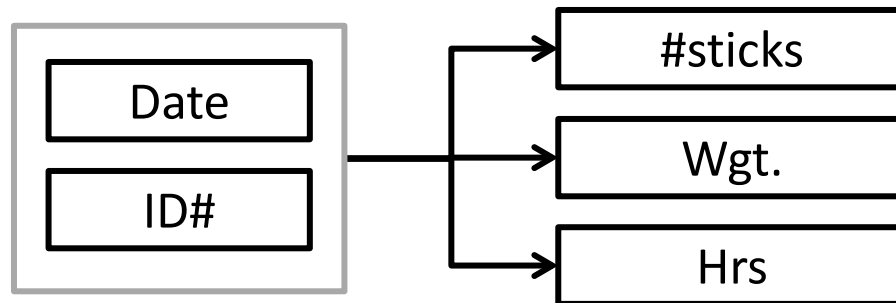




# Normal Forms

## 2<sup>nd</sup> Normal Form (2NF)

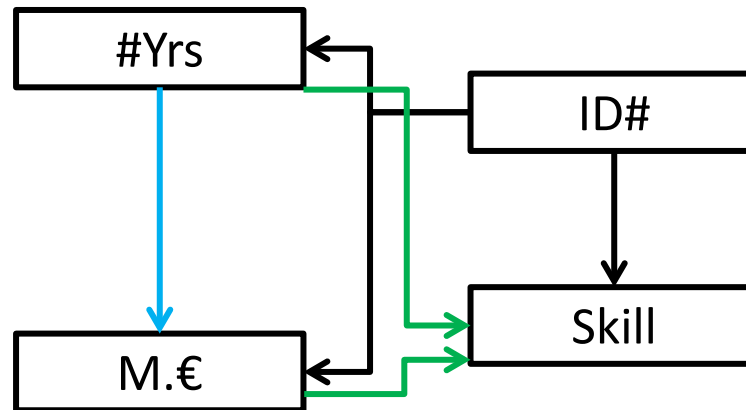
- A relation is in 2NF *if and only if* every nonkey attribute is *irreducibly dependent* on the Primary Key



# Normal Forms

- What's wrong with this relation?

ID#	Skill	#Yrs	M.€
1	Medium	2	40
2	Low	1	30
3	High	3	45
4	High	4	50

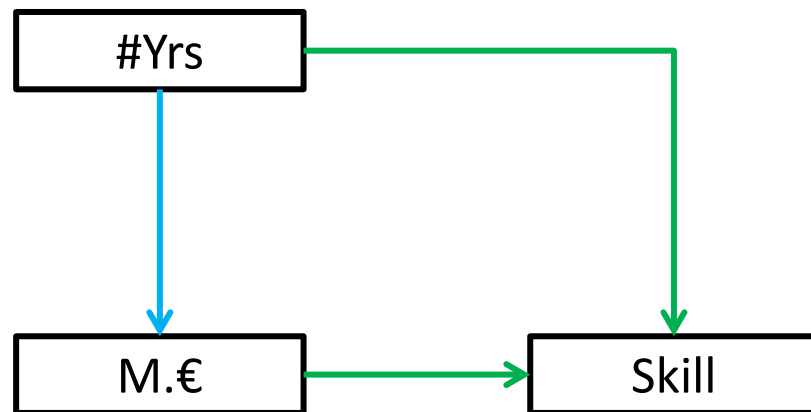
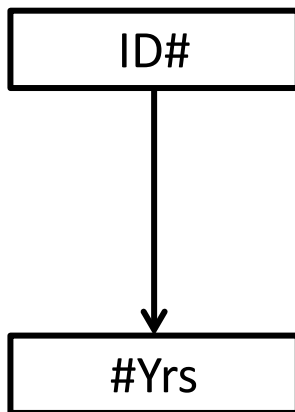


# Normal Forms

- Decompose the relation – again by projection

ID#	#Yrs
1	2
2	1
3	2
4	4

#Yrs	Skill	M.€
2	Medium	40
1	Low	30
3	High	45
4	High	50



# Normal Forms

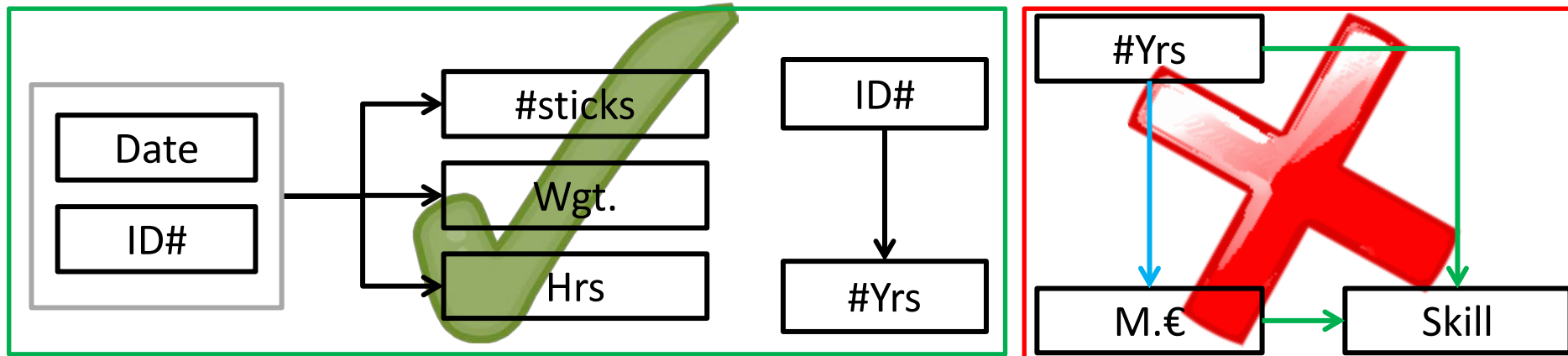
## 3<sup>rd</sup> Normal Form (3NF)

- A relation is in 3NF ***if and only if*** it is in 2NF every nonkey attribute is ***nontransitively dependent*** on the Primary Key

# Normal Forms

## 3<sup>rd</sup> Normal Form (3NF)

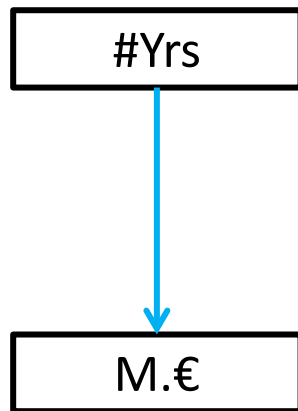
- A relation is in 2NF ***if and only if*** it is in 2NF every nonkey attribute is ***nontransitively dependent*** on the Primary Key



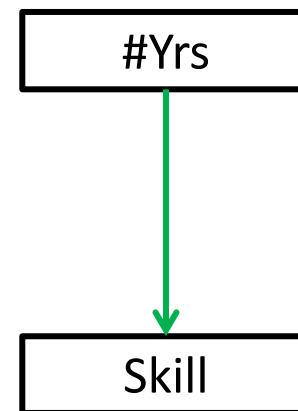
# Normal Forms

- Decompose the relation – again by projection

#Yrs	M.€
2	40
1	30
3	45
4	50



#Yrs	Skill
2	Medium
1	Low
3	High
4	High



# Boyce-Codd Normal Form

- Note in the previous examples we considered only a single candidate key
- Boyce-codd normal form considers also cases where we have overlapping candidate keys

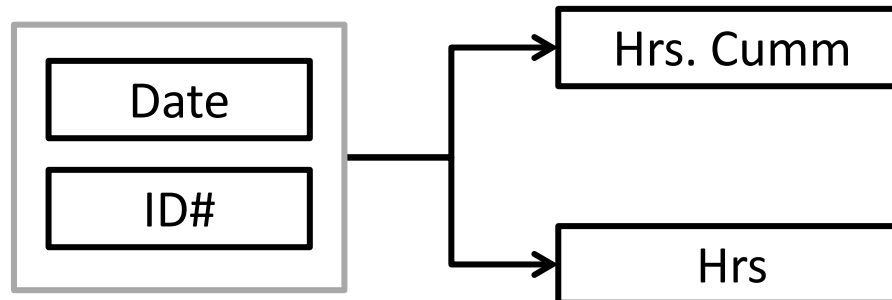
## Boyce-Codd Normal Form (BCNF)

- A relation is in BCNF *if and only if* every **nontrivial left irreducible** FD has a candidate key as its determinant (LHS)

# Boyce-Codd Normal Form

- In a diagram

ID#	Date	#sticks	Wgt.	Hrs	Hrs. Cumm
1	1.06	55	9	6	2212
2	7.05	34	5	5	3182
3	1.06	54	9	6	3097
4	3.11	61	12	8	5220

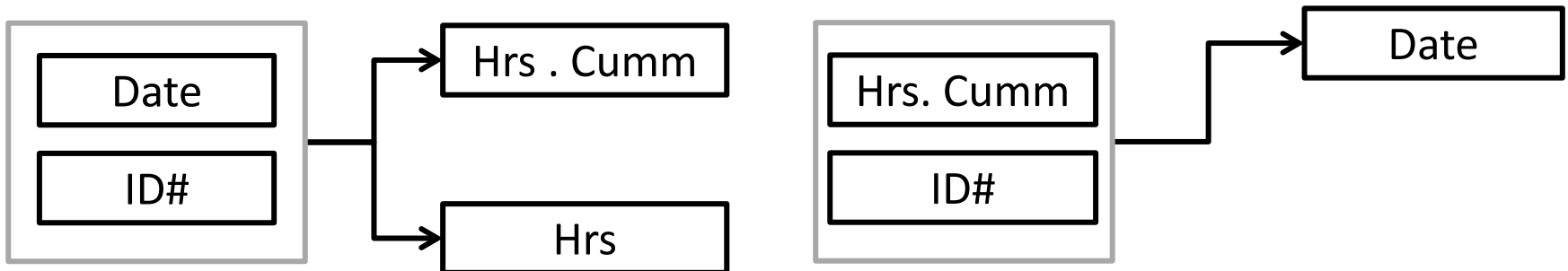




# Boyce-Codd Normal Form

- In a diagram

ID#	Date	#sticks	Wgt.	Hrs	Hrs. Cumm
1	1.06	55	9	6	2212
2	7.05	34	5	5	3182
3	1.06	54	9	6	3097
4	3.11	61	12	8	5220



# References

- C.J. Date, *An Introduction to Database Systems, 8<sup>th</sup> Edition*. Pearson Education Inc., 2004.
- See [www.geoinformatic.cc](http://www.geoinformatic.cc)

That's **NOT** all for today

Practical

That's all for today

Thank you!

Questions?